

Answers to HW Set #2

2.1. Buying fruit.

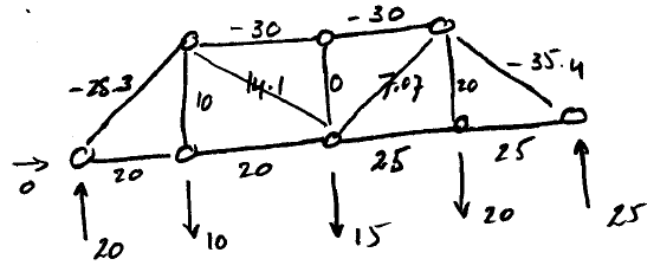
```
A = [3 12 1; 12 0 2; 0 2 3]
b = [2.36; 5.26; 2.77]
format bank
x = A\b
x =
0.29
0.05
0.89
```

2.2. Reduced row echelon form of the magic square of order six.

```
rref(magic(6))
1 0 0 0 0 -2
0 1 0 0 0 -2
0 0 1 0 0 1
0 0 0 1 0 2
0 0 0 0 1 2
0 0 0 0 0 0
```

2.3 I would add unknown support forces at ① and ② plus node equilibrium at ③ and ④

① $H_1 + \alpha f_1 + f_2 = 0$
 $V_1 + \alpha f_1 = 0$
 ③ $V_8 + \alpha f_2 = 0$



2.5 a) magic(n) NO
 hilb(n) YES
 pascal(n) YES
 eye(n,n) YES
 rand(n,n) NO

b) note:

 $R^T * R = A \rightarrow$

$R = \text{randn}(n,n); A = R * R$ YES
 $R = \text{randu}(n,k); A = R' * R$ MAYBE
 $R = \text{randu}(n,n); A = R' * R + \text{eye}(n,n)$ MAYBE

```
function R = mychol(A)
%MYCHOL My own version of Cholesky factorization.
% R = mychol(A), for a positive definite matrix A,
% is an upper triangular matrix R so that R'*R = A.

if ~isequal(A',A)
    error('Matrix must be symmetric or Hermitian')
end
[n,n] = size(A);
R = zeros(n,n);
for k = 1:n
    i = (1:k-1)';
    j = k+1:n;
    s = A(k,k) - R(i,k)'*R(i,k);
    if s > 0
        R(k,k) = sqrt(s);
    else
        error('Matrix must be positive definite')
    end
    R(k,j) = (A(k,j) - R(i,k)'*R(i,j))/R(k,k);
end
```

2.9 a) $\underline{x} = \underline{p} + t \cdot \underline{z}$ with $A\underline{p} = \underline{b}$ and $A\underline{z} = \underline{0}$
 particular sol $\underline{p} = (1/3 \ 1/3 \ 0)^T$ and null space $\underline{z} = (1 \ -2 \ 1)^T$
 b) finally you end up at $v(3,3) = 0$ and $c(3) = 0 \rightarrow$
 $x(3) = 0$ an arbitrary value, which is the free parameter.
 c) $\underline{x} = (13/13 \ -25/13 \ 4)^T$ Good, because residual $r = b - A * x$ is small
 Bad because sol is not unique!
 Due to finite precision we end up with $x(3) = \frac{2 * \text{eps}}{\text{eps}/2} = 4$
 d) \ uses column-oriented algorithm $\rightarrow x(3) = \frac{54 * \text{eps}}{\text{eps}/2} = 5/2!$
 bslashtx uses row-oriented algorithm

2.19 a) $n = 100; e = \text{ones}(n,1); b = (1:n)';$
 $A = 2 * \text{diag}(e) - \text{diag}(e(1:n-1), -1) - \text{diag}(e(1:n-1), 1);$
 $x_a = \text{bslash}x(A,b);$

b) $A = \text{spdiags}([-e \ 2 * e \ -e], [-1 \ 0 \ 1], n, n);$
 $x_b = A \backslash b;$

c) $x_c = \text{tridissolve}(-e, 2 * e, -e, b);$

d) $\text{condost}(A) \rightarrow 5.1 \cdot 10^3$

2.21 $U\{k\}$ is a string, the k^{th} URL.
 $U(k)$ is a 1-by-1 cell array containing the string of the k^{th} URL.
 $G(k,:)$ is nonzero for incoming nodes to the k^{th} URL.
 $G(:,k)$ is nonzero for outgoing nodes from the k^{th} URL.
 $U(G(k,:))$ is a list of incoming URL's to the k^{th} URL*.
 $U(G(:,k))$ is a list of outgoing URL's from the k^{th} URL.

*Note that this only works if G is logical.

2.22. Cliques in the harvard500 Web connectivity matrix.
 $U(168:180)$: Harvard Divinity School
 $U(229:248)$: Radcliffe Institute
 $U(261:281)$: Dana-Farber Cancer Institute
 $U(315:335)$: "Go Crimson", Harvard's athletic program

2.23. (a) For $p \geq 8$, $\text{nnz}(G^p) = 167985$.
 (b) $\text{nnz}(G^8)/\text{prod}(\text{size}(G)) = 0.6719$.
 (c) for $p = 1:9$, `subplot(3,3,p), spy(G^p), end`
 (d) The "Go Crimson" athletic program, nodes 46 and 315:335, has no links to the other pages in the data set.

2.25. Disconnected miniweb.

```
G =
  0 0 0 1 0 0
  1 0 0 0 0 0
  1 1 0 0 0 0
  0 1 1 0 0 0
  0 0 0 0 0 1
  0 0 0 0 1 0
pagerank1(G,.85) =
  0.1981
  0.1092
  0.1556
  0.2037
  0.1667
  0.1667
```

What happens to page rank as $p \rightarrow 1$? Two possible answers here. The intuitive answer is that the graph has two disconnected subgraphs and consequently the Markov stationary probabilities are not unique. The direct solution algorithm used in `pagerank` certainly breaks down if $p = 1$. However, a second answer is that `pageranksym`, a symbolic version of `pagerank1`, produces

```
p = sym('p');
pagerank1(G,p) =
 [ 1/3*(p^3+3*p^2+2*p+2)/(p^3+4*p^2+4*p+4) ]
 [ 1/3*(p^2+p+2)/(p^3+4*p^2+4*p+4) ]
 [ 1/6*(p^3+3*p^2+4*p+4)/(p^3+4*p^2+4*p+4) ]
 [ 1/6*(p^3+5*p^2+6*p+4)/(p^3+4*p^2+4*p+4) ]
 [ 1/6 ]
 [ 1/6 ]
```

```
and
limit(ans,p,1) =
 [ 8/39 ]
 [ 4/39 ]
 [ 6/39 ]
 [ 8/39 ]
 [ 1/6 ]
 [ 1/6 ]
```

These values are $2/3$ times the limiting values for the 4-by-4 subgraph and $1/3$ times the values for the 2-by-2 subgraph.